## Design and Analysis of Algorithms

## Unit - I



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## Algorithm and Analysis

## Syllabus <br> UNIT -I: ALGORITHM AND ANALYSIS

What is an Algorithm? - Algorithm Specification - Performance Analysis - Randomized Algorithms.

## TEXT BOOK

Fundamentals of Computer Algorithms, Ellis Horowitz, Sartaj Sahni, Sanguthevar Rajasekaran, Galgotia Publications, 2015.

## Introduction to the Concept of Algorithms

- Algorithm
- Problem Solving
- Design of an Algorithm
- Analysis of an algorithm


## Notion of an Algorithm



## Algorithm

- An algorithm is a finite set of instructions that, if followed, accomplishes a particular task i.e., for obtaining a required output for any legitimate input in a finite amount of time.
- All algorithms must satisfy the following criteria:
- Definiteness. Each instruction is clear and unambiguous.
- Effectiveness. Every instruction must be very basic so that it can carried out, by a person using pencil and paper.
- Finiteness. If we trace out the instructions of an algorithm, then for all cases, the algorithm terminates after a finite number of steps.
- Input. Zero or more quantities are externally supplied.
- Output. At least one quantity is produced.


## Algorithm Specification

- An algorithm can be described in three ways:
- Natural language in English
- Graphic representation called flowchart
- Pseudo-code method
$>$ In this method we typically represent algorithms as program, which resembles C language

1. Input two numbers
2. Add the two numbers
3. Print the result


## Pseudo-code Conventions

1. Comments begin with // and continue until the end of line.
2. Blocks are indicated with matching braces $\{$ and $\}$.
3. An identifier begins with a letter. The data types of variables are not explicitly declared.
4. Assignment of values to variables is done using the assignment statement.

〈variable〉:= «expression»;
5. There are two Boolean values true and false.
$>$ Logical operators: AND, OR, NOT
$>$ Relational operators: $<, \leq,=, \neq,>, \geq$

## Pseudo－code Conventions

6．The following looping statements are used： while，for and repeat－until

repeat－until：
repeat
〈statement 1〉
。
－
〈statement n〉
until «condition»

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## Pseudo－code Conventions

7．A conditional statement has the following forms：
if＜condition» then 〈statement〉
if «condition» then «statement 1 〉 else＜statement 2）
case statement：
case
\｛
：＜condition 1»：＜statement 1＞
：＜condition n ：〈statement n 〉
：else：〈statement $\mathrm{n}+1$ 〉
\}

## Pseudo-code Conventions

8. Input and output are done using the instructions read and write.
9. There is only one type of procedure: Algorithm.

Algorithm contains
$>$ Heading
$>$ Body
The heading takes the form
Algorithm Name («parameter list») $\longrightarrow$ heading \{

$\}$

## Pseudo-code Conventions

1. Algorithm $\operatorname{Max}(\mathrm{A}, \mathrm{n})$
2. // A is an array of size n .
3. \{
4. Result $:=\mathrm{A}[1]$;
5. for $\mathrm{i}:=2$ to n do
6. if $\mathrm{A}[\mathrm{i}]>$ result then
7. Result := A[i];
8. return Result;
9. \}

## Performance Analysis

1. Space Complexity
2. Time Complexity

Space complexity of an algorithm is the amount of memory it needs to run to complete.
Space needed by an algorithm is given by $\mathrm{S}(\mathrm{P})=\mathrm{C}($ fixed part $)+\mathrm{Sp}($ variable part $)$
fixed part: independent of instance characteristics. Eg. Space for simple variables, constants etc.
variable part: space for variables whose size is dependent on particular problem instance

## Performance Analysis



```
Algorithm-2
Algorithm Sum(a, n)
{
s:=0.0;
for i:=1 to n do
    s:= s + a[i];
return s;
}
```


## Performance Analysis

| ```Algorithm-3 Algorithm RSum(a, n) \{ if \((\mathrm{n} \leq 0)\) then return 0.0; else Return RSum(a, n-1)+a[n]; \}``` |  |
| :---: | :---: |
|  |  |
|  | $\operatorname{RSum}(\mathrm{a}, \mathrm{n}-1)=1(\mathrm{a}[\mathrm{n}-1])+1(\mathrm{n})+1($ return $)$ |
|  | ........... |
|  | $\ldots \ldots \ldots \ldots$. $\operatorname{RSum}(\mathrm{a}, \mathrm{n}-\mathrm{n})=1(\mathrm{a}[\mathrm{n}-\mathrm{n}])+1(\mathrm{n})+1($ return $)$ |
|  | Total $\rightarrow \geq 3$ (n+1) units |

## Performance Analysis

## 2. Time Complexity

The time complexity of an algorithm is the amount of computer time it needs to run to complete.
$\mathrm{T}(\mathrm{P})=$ compile time + execution time
$\mathrm{T}(\mathrm{P})=\mathrm{Tp}($ execution time $)$

## Step count:

$>$ For algorithm heading $\rightarrow 0$
$>$ For braces $\rightarrow 0$
$>$ For expressions $\rightarrow 1$
$>$ For any looping statements $\rightarrow$ number of times the loop is repeating

## Performance Analysis

## Algorithm-1

Algorithm abc(a,b,c)
\{
return $\mathrm{a}+\mathrm{b}+\mathrm{b}^{*} \mathrm{c}+(\mathrm{a}+\mathrm{b}-\mathrm{c}) /(\mathrm{a}+\mathrm{b})+4.0$; \}

$$
\left[\begin{array}{l}
\rightarrow 0 \\
\rightarrow 0 \\
\rightarrow 1 \\
\rightarrow 0 \\
1 \text { unit }
\end{array}\right.
$$

| Algorithm-2 |  |
| :---: | :---: |
| Algorithm Sum(a, n) | $\rightarrow 0$ |
| \{ | $\rightarrow 0$ |
| s: $=0.0$; | $\rightarrow 1$ |
| for $\mathrm{i}:=1$ to n do | $\zeta \rightarrow \mathrm{n}+1$ |
| $\mathrm{s}:=\mathrm{s}+\mathrm{a}[\mathrm{i}]$; | n |
| return s; | $\rightarrow 1$ |
| \} | $\rightarrow 0$ |
| Dr. R. Bhuvaneswari | $2 \mathrm{n}+3$ units |

## Performance Analysis

## Algorithm-3 <br> Algorithm RSum(a, n) <br> \{ <br> if $(\mathrm{n} \leq 0)$ then <br> return 0.0; <br> else <br> return $\operatorname{RSum}(\mathrm{a}, \mathrm{n}-1)+\mathrm{a}[\mathrm{n}]$; \}

$$
\begin{array}{rlr}
\mathrm{T}(\mathrm{n})= & \text { if } \mathrm{n}=0 \\
& =2+\mathrm{T}(\mathrm{n}-1) \quad \text { if } \mathrm{n}>0 \\
\mathrm{~T}(\mathrm{n})= & 2+\mathrm{T}(\mathrm{n}-1) \\
= & 2+(2+\mathrm{T}(\mathrm{n}-2)) \\
= & 2+2+T(\mathrm{n}-2)=2 * 2+\mathrm{T}(\mathrm{n}-2) \\
= & 2 * 2+(2+T(\mathrm{n}-3)) \\
= & 2 * 2+2+T(\mathrm{n}-3)=2 * 3+\mathrm{T}(\mathrm{n}-3) \\
& \cdots \cdots \cdots \cdots \\
& \cdots \cdots \cdots \cdots \\
= & 2 * \mathrm{n}+\mathrm{T}(\mathrm{n}-\mathrm{n})=2 \mathrm{n}+\mathrm{T}(0) \\
\mathrm{T}(\mathrm{n})= & 2 \mathrm{n}+2 \text { units }
\end{array}
$$

## Randomized algorithms

- Makes use of randomizer (random number generator).
- Decisions made in the algorithm depends on the output of the randomizer.
- Output and execution time may very from run to run for the same input.


## Randomized algorithms

\{
while(true) do
\{
$\mathrm{i}=\operatorname{Random}() \bmod \mathrm{n}+1$;
$\mathrm{j}=$ Random( $) \bmod \mathrm{n}+1$;
if $((\mathrm{i} \# \mathrm{j})$ and $(\mathrm{a}[\mathrm{i}]=\mathrm{a}[\mathrm{j}]))$ then
return i;
\}
\}

Algorithm RepeatedElement(a,n)
$\mathrm{i}=\mathbf{1 , j} \mathbf{j}=\mathbf{6}$
1 \# 6 and a[1] \# a[6]
$\mathrm{i}=\mathbf{1 , j} \mathbf{j}=5$
1 \# 5 and a[1] \# a[5]
$\mathrm{i}=\mathbf{2}, \mathrm{j}=2$
$2=2$
$\mathrm{i}=4, \mathrm{j}=9$
4 \# 9 and a[4] \# a[9]
$\mathbf{i}=\mathbf{9}, \mathbf{j}=\mathbf{3}$
9 \# 3 and a[9] \# a[3]
$\mathrm{i}=\mathbf{6}, \mathrm{j}=\mathbf{7}$
6 \# 7 and $\mathrm{a}[6]=\mathrm{a}[7]$
$a[1] \quad a[2] \quad a[3] \quad a[4] \quad a[5] \quad a[6] \quad a[7] \quad a[8] \quad a[9] \quad a[10]$

| 10 | 20 | 30 | 40 | 50 | 60 | 60 | 60 | 60 | 60 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

Eg.

